

# MATH 301

## INTRODUCTION TO PROOFS

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- Relations
- Functions

## Relevant sections of the textbook

- Chapter 3
- Chapter 5

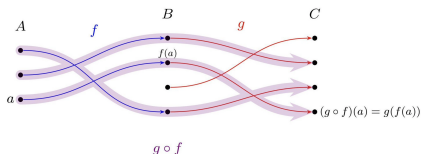
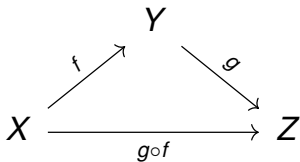
## Recall: Compositionality of functions

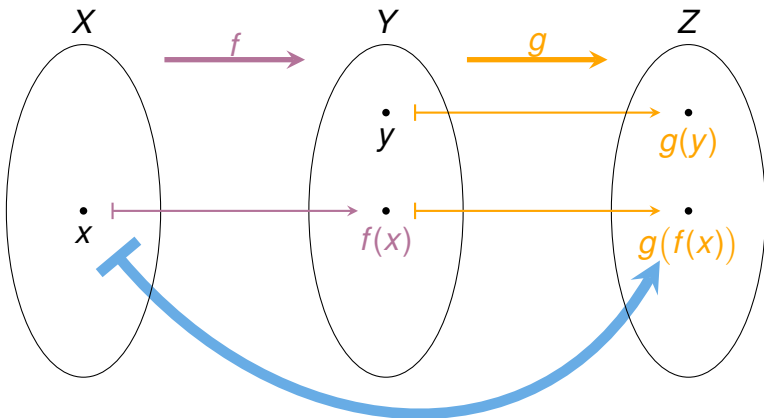
For any set  $X$ , we can define a function  $\text{id}: X \rightarrow X$  by letting  $\text{id}(x)$  to be the same as  $x$ . This function is called the **identity** function on  $X$ .

More interestingly, let  $f: X \rightarrow Y$  and  $g: Y \rightarrow Z$  be functions. We can define a new function  $k: X \rightarrow Z$  by letting

$$k(x) =_{\text{def}} g(f(x))$$

The function  $k$  is called the **composition** of  $f$  and  $g$  which we also call “ $f$  composed with  $g$ ” (or “ $g$  after  $f$ ”) and which we denote by  $g \circ f$ .





$$\lambda y. g(y) \circ \lambda x. f(x) = \lambda x. g[f(x)/y]$$

$$\lambda y. \log_2 y \circ \lambda x. 2^x = \lambda x. \log_2 y [2^x/y] = \log_2 2^x = x$$

The composition of function introduced above has two important properties:

**unitality** for any function  $f: X \rightarrow Y$ , we have  $f \circ \text{id}_X = f$  and  $\text{id}_Y \circ f = f$ .

**associativity** for any functions  $f: W \rightarrow X$ ,  $g: X \rightarrow Y$  and  $h: Y \rightarrow Z$ , we have

$$h \circ (g \circ f) = (h \circ g) \circ f.$$

For any function  $f: X \rightarrow Y$ , we define as subset of  $X \times Y$  known as the **graph** of  $f$ .

$$\mathbf{Gr}(f) = \{(x, y) \mid f(x) = y\}$$

Define functions  $h$ ,  $i$ , and  $p$  as follows:

$$h = \lambda x.(x, f(x)) \quad (1)$$

$$i = \lambda(x, y).(x, y) \quad (2)$$

$$p = \lambda(x, y).y \quad (3)$$

## Exercise

Show that the functions  $f$ ,  $h$ ,  $i$ , and  $p$  fit into the following square of sets and functions commutes:

$$\begin{array}{ccc} \mathbf{Gr}(f) & \xrightarrow{i} & X \times Y \\ h \uparrow & & \downarrow p \\ X & \xrightarrow{f} & Y \end{array}$$

## Composition of relations

Given a relation  $R$  on  $X$  and  $Y$  and a relation  $S$  on  $Y$  and  $Z$  we can compose them to get a relation  $S \circ R$  on  $X$  and  $Z$  defined as follows:

$$x(S \circ R)z \iff \exists y \in Y (xRy \wedge yRz)$$

### Exercise

Let  $B$  be the “brotherhood” relation ( $xBy$  means  $x$  is a brother of  $y$ ) and  $S$  be the “sistership” relation. Show that the composite relation  $S \circ B$  is not equivalent to  $B$ .

### Exercise

- Prove that if both  $R$  and  $S$  are partial orders then  $S \circ R$  is a partial order.
- Prove that if both  $R$  and  $S$  are equivalence relations then  $S \circ R$  is an equivalence relation.

## Exercise

Show that for any equivalence relation  $R$  on a set  $X$  we have

①  $R \circ R = R.$

②  $R \circ R \circ \dots \circ R = R$



# Composition of functions from compositions of relations

## Theorem

*Suppose  $f: X \rightarrow Y$  and  $g: Y \rightarrow Z$  are functions. Consider the corresponding relations  $R_f$  and  $R_g$ . The relation corresponding to the composite function  $g \circ f$  is equivalent to the composite relations  $R_g \circ R_f$ , that is,*

$$\forall x \in X \forall z \in Z (x R_{g \circ f} z \iff x (R_g \circ R_f) z)$$

# Isomorphisms of sets

## Definition

An *isomorphism* between two sets  $X$  and  $Y$  is a pair of function

$$f: X \rightarrow Y \text{ and } g: Y \rightarrow X$$

such that  $g \circ f = \text{id}_X$ , and  $f \circ g = \text{id}_Y$ .

We can think of functions  $f$  and  $g$  above as no data-loss “processes”, e.g. conversion of files to different format without data being lost.

## Definition

The sets  $X$  and  $Y$  are said to be *isomorphic* in case there exists an isomorphism between them. In this case, we use the notation  $X \cong Y$ .

## Exercise

*Show that for any set  $A$ , it is isomorphic to  $\emptyset$  if and only if  $A$  does not have any elements. Can you prove this without the LEM?*

Previously, we defined the cartesian product  $A \times B$  of two sets  $A$  and  $B$  to consists of all the pairs  $(a, b)$  where  $a \in A$  and  $b \in B$ . Now, we show that if we have more two sets the order of forming products does not matter.

### Exercise

- 1 For all sets  $A, B, C$  we have

$$(A \times B) \times C \cong (A \times B) \times C$$

For this reason, we use  $A \times B \times C$  to denote either sets.

## Exercise

*Show that two finite sets are isomorphic if and only if they have the same number of elements.*

## Exercise

Show that for any function  $f: X \rightarrow Y$ , we have

$$\mathbf{Gr}(f) \cong X.$$

## A remark on disjoint unions

We introduced the operation of taking disjoint union of two sets as follows:

$$A \sqcup B = \{\text{inl}(x) \mid x \in A\} \cup \{\text{inr}(x) \mid x \in B\}$$

### Exercise

Show that

$$A \sqcup B \cong (\{0\} \times A) \cup (\{1\} \times B)$$

Inspired by this fact we define the **disjoint union of a family**  $\{A_i \mid i \in I\}$  of sets to be

$$\bigsqcup_{i \in I} A_i = \bigcup_{i \in I} \{i\} \times A_i.$$

An element of  $\bigsqcup_{i \in I} A_i$  is a pair  $(i, a)$  where  $i \in I$  and  $a \in A_i$ .

## Inverse of a relation

We can always define an inverse to a relation:

### Definition

*For a relation  $R$  on  $X$  and  $Y$  we define the inverse of  $R$  to be a relation  $R^{-1}$  on  $Y$  and  $X$  defined by*

$$yR^{-1}x \Leftrightarrow xRy$$

### Exercise

*Show that if a relation  $R$  is functional then it is not necessarily the case that  $R^{-1}$  is functional.*



## Arithmetic of sets

We define the operation of addition on sets as follows: For sets  $X$  and  $Y$  let the sum  $X + Y$  be defined by their disjoint union  $X \sqcup Y$ .

### Exercise

- 1 Show that the addition operation on sets is both commutative and associative.
- 2 Show that the empty set is the unit (aka neutral element) of addition of sets.

### Exercise

Show that  $\underline{m} + \underline{n} \cong \underline{m + n}$  for all natural numbers  $m$  and  $n$ .

## Exercise

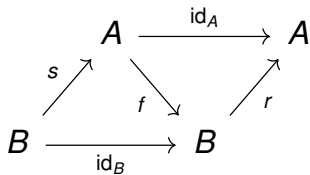
- 1 Show that if  $S$  and  $S'$  are isomorphic, then for all sets  $X$ , we have  $X + S \cong X + S'$ .
- 2 Prove that for any singleton  $S$ , we have  $\mathbb{N} + S \cong \mathbb{N}$ .

Sometimes, when the context precludes risk of confusion, we use the notation  $1$  for any singleton set. Therefore, we can simplify the last statement in above to

$$\mathbb{N} + 1 \cong \mathbb{N}.$$

## Definition

- A **retract** (aka **left inverse**) of function  $f: A \rightarrow B$  is a morphism  $r: B \rightarrow A$  such that  $r \circ f = \text{id}_A$ . In this case we also say  $A$  is a retract of  $B$ .
- A **section** (aka **right inverse**) of function  $f: A \rightarrow B$  is a morphism  $s: B \rightarrow A$  such that  $f \circ s = \text{id}_B$ .



## Example

- The circle is a retract of punctured disk.
- The maps from the infinite helix to the circle has a section, but no continuous section.

# Injections

## Definition

A function  $f: X \rightarrow Y$  is *injective* (or *one-to-one*) if

$$\forall a, b \in X (f(a) = f(b) \Rightarrow a = b)$$

An injective function is said to be an *injection*.

# Surjections

## Definition

A function  $f: X \rightarrow Y$  is *surjective* (aka *onto*) if

$$\forall y \in Y, \exists x \in X, f(x) = y$$

holds. A surjective function is said to be a *surjection*.

## Proposition

- ① *A function with a retract is injective.*
- ② *A function with a section is surjective.*

## Injection and retracts

Does every injection have a retract?

## Injection and retracts

No. Consider the function  $\emptyset \rightarrow \mathbf{1}$ .



## Injection and retracts

### Proposition

*Let  $f : X \rightarrow Y$  be a function. If  $f$  is injective and  $X$  is inhabited, then  $f$  has a retract.*

## Injection and retracts

### Proof.

Suppose that  $f$  is injective and  $X$  is inhabited. Since  $X$  is inhabited, we get always fix an element of it, say  $x_0 \in X$ . Now, define  $r: Y \rightarrow X$  as follows.

$$r(y) = \begin{cases} x & \text{if } y = f(x) \text{ for some } x \in X \\ x_0 & \text{otherwise} \end{cases}$$

Note that  $r$  is well-defined since if for some  $y$ , there are elements  $x$  and  $x'$  such that  $y = f(x) = f(x')$ , then, by injectivity of  $f$ , we have  $x = x'$ , and therefore, the value of  $r$  is uniquely determined.

To see that  $r$  is a retract of  $f$ , let  $x \in X$ . Letting  $y = f(x)$ , we see that  $y$  falls into the first case in the specification of  $r$ , so that  $r(f(x)) = g(y) = a$  for some  $a \in X$  for which  $y = f(a)$ . But,  $f(x) = y = f(a)$ , and by injectivity of  $f$  we have  $x = a$ . Therefore, for every  $x \in X$ ,

## Injection and retracts

Was this proof constructive?

# Questions

Time for your questions!