

MATH 301

INTRODUCTION TO PROOFS

Sina Hazratpour

Johns Hopkins University

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- Recursion

Relevant sections of the textbook

- Chapter 4

Overview

- 1 Recursion
- 2 Applications of recursion theorem
- 3 Recursion in Lean

Recall that in the last lecture, we defined the set of **natural numbers** \mathbb{N} to be a set *generated* by the by the number 0 and the **successor function** $\text{succ}: \mathbb{N} \rightarrow \mathbb{N}$.

This means that the only way to construct a natural number is to start with 0 and apply the successor function finitely many times. Therefore, natural numbers are

$$0, \text{succ}(0), \text{succ}(\text{succ}(0)), \dots$$

We also postulated the principle of induction on natural numbers.

Predicates and subsets

Recall that we can think of a predicate P on natural numbers as a function $P: \mathbb{N} \rightarrow \mathbf{2}$ where the set $\mathbf{2}$ consists of truth values \perp and \top .

Note that the set $\mathbb{N} \rightarrow \mathbf{2}$ is in bijection with $\mathcal{P}(\mathbb{N})$.

In one way, we construct a function

$$\eta: (\mathbb{N} \rightarrow \mathbf{2}) \rightarrow \mathcal{P}(\mathbb{N})$$

whose value at a predicate P is the set consisting of all $n \in \mathbb{N}$ such that $P(n)$ is true, i.e.

$$\eta(P) =_{\text{def}} \{n \in \mathbb{N} \mid P(n)\}$$

In the other direction, we take a subset S of \mathbb{N} to the characteristic function $\chi_S: \mathbb{N} \rightarrow \mathbf{2}$.

The principle of induction

The principle of induction says that for any property $P: \mathbb{N} \rightarrow \mathbf{2}$ of natural numbers, if

- 1 $P(0)$ holds, and
 - 2 whenever $P(n)$ holds then $P(n + 1)$ holds,
- we have that P holds of every natural number.

The principle of induction reformulated

Let $S \subseteq \mathbb{N}$ be any set of natural numbers that contains 0 and is closed under the successor operation. Then $S = \mathbb{N}$.

Proofs vs computation

We saw that the principle of induction is a very powerful tool in **proving** universally quantified statements about natural numbers.

Example

- *For any finite set S , if S has n elements, then there are 2^n subsets of S .*
- *For every $n \in \mathbb{N}$, we have $0^2 + 1^2 + 2^2 + \dots + n^2 = \frac{1}{6} n(n+1)(2n+1)$.*

But, we also need to **compute** with natural numbers. At the very least, we should be able to define the arithmetic operations $+$, \times , etc.

That is why we need another principle to help us with computation of natural numbers. This is the so-called principle of **recursion** which in fact can be proved from the principle of induction!

Recursion theorem

Theorem

Let A be a set. For all $a \in A$ and all $g: \mathbb{N} \times A \rightarrow A$, there is a unique function $f: \mathbb{N} \rightarrow A$ such that

- 1 $f(0) = a$
- 2 $f(\text{succ}(n)) = g(n, f(n))$ for all $n \in \mathbb{N}$.

Proof.

Theorem 4.1.2 (Recursion theorem) Page 145. □

Since for every function g such function f is uniquely determined, we write $\text{rec}(g)$ for it.

We have

$$\text{rec}(g)(0) = a$$

$$\text{rec}(g)(1) = \text{rec}(g)(\text{succ}(0)) = g(0, \text{rec}(g)(0)) = g(0, a)$$

$$\text{rec}(g)(2) = \text{rec}(g)(\text{succ}(1)) = g(1, \text{rec}(g)(1)) = g(1, g(0, a))$$

Recursion, practically!

In order to specify a function $f: \mathbb{N} \rightarrow A$, it suffices to define $f(0)$ and, for given $n \in \mathbb{N}$, assume that $f(n)$ has been defined, and define $f(\text{succ}(n))$ in terms of n and $f(n)$.

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Addition by recursion

We define additions of natural numbers as a function $+: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$.

This means for every $m \in \mathbb{N}$, we have to define a function $m + (-): \mathbb{N} \rightarrow \mathbb{N}$.

We define the latter by recursion: consider the function $g: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ given by the assignment $g(i, j) = \text{succ}(j)$.

Choose $a = m$ in the recursion theorem. Therefore,

$$\text{rec}(g)(0) = m$$

$$\text{rec}(g)(1) = \text{rec}(g)(\text{succ}(0)) = g(0, \text{rec}(g)(0)) = g(0, m) = \text{succ}(m)$$

$$\text{rec}(g)(2) = \text{rec}(g)(\text{succ}(1)) = g(1, \text{rec}(g)(1)) = \text{succ}(\text{succ}(m))$$

$$\text{rec}(g)(\text{succ}(n)) = g(n, \text{rec}(g)(n))$$

Addition by recursion

We now define $m + (-) : \mathbb{N} \rightarrow \mathbb{N}$ to be $\text{rec}(g) : \mathbb{N} \rightarrow \mathbb{N}$.

$$m + 0 = m \tag{1}$$

$$m + \text{succ}(n) = \text{succ}(m + n) \tag{2}$$

Therefore,

$$m + 1 = m + \text{succ}(0) = \text{succ}(m + 0) = \text{succ}(m) \tag{3}$$

In particular,

$$1 + 1 = \text{succ}(1) = \text{succ}(\text{succ}(0)) = 2$$

Combining recursion and induction

Proposition

For every natural numbers m , we have $m + 1 = 1 + m$.

Proof.

We use induction on m to prove that $m + 1 = 1 + m$ for all $m \in \mathbb{N}$.

When $m = 0$, by equations (1) and (2), we have

$$1 + 0 = 1 = \text{succ}(0) = \text{succ}(0 + 0) = 0 + \text{succ}(0) = 0 + 1.$$

Suppose that $1 + m = m + 1$. We want to show that $1 + \text{succ}(m) = \text{succ}(m) + 1$.

But, by definition of function $m + (-)$ for $m = 1$,

$$1 + \text{succ}(m) = \text{succ}(1 + m) = \text{succ}(m + 1) = \text{succ}(\text{succ}(m)) = \text{succ}(m) + 1,$$

where the last two equations above follow from equation (3). □

Proposition (commutativity of addition of natural numbers)

For every natural numbers m and n , we have $m + n = n + m$.

Proof left to the reader.

Hint: We prove, by induction, the following lemmas first:

Lemma (neutrality of 0 for +)

For all natural numbers k we have $k + 0 = 0 + k$.

Lemma (associativity of addition)

For all natural numbers $k + (m + n) = (k + m) + n$.

Proof.

We prove the commutativity of addition by fixing m and inducting on n .

If $n = 0$, by the neutrality of 0 (lemma above) we have that $m + 0 = 0 + m$, and we are done. Suppose that $m + n = n + m$. We want to prove that $m + \text{succ}(n) = \text{succ}(n) + m$.

$$\begin{aligned} m + \text{succ}(n) &= m + (n + 1) && \text{by eq (3)} \\ &= (m + n) + 1 && \text{by associativity of addition} \\ &= (n + m) + 1 && \text{by inductive hypothesis} \\ &= n + (m + 1) && \text{by associativity of addition} \\ &= n + (1 + m) && \text{by the last proposition} \\ &= (n + 1) + m && \text{by associativity of addition} \\ &= \text{succ}(n) + m && \text{by eq (2)} \end{aligned}$$

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